

Tropos: Goal Analysis



Paolo Giorgini

Department of Information and
Communication Technology

University of Trento - Italy

<http://www.dit.unitn.it/~pgiorgio>

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Outline

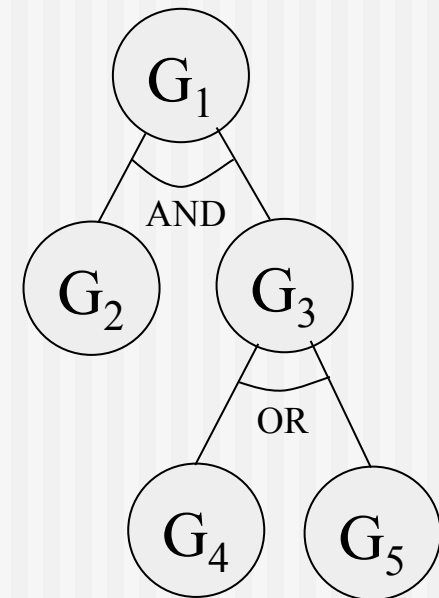
- The concept of goal in SE
 - Goal analysis
- Tropos Goal Reasoning Framework
 - Qualitative approach
 - Quantitative approach
 - Backward and forward reasoning
 - GR-Tool

Goals

- The concept of goal has been used in many area of Computer Science, e.g.,
 - in planning to describe desirable states of the world
 - in agent architecture to describe agents' mental state
- More recently, goals have been used in Software Engineering to model:
 - Early requirements (e.g., *every book request will eventually be fulfilled*)
 - Non-functional requirements (e.g., *the new system will be highly reliable*)

Goal Analysis

- Traditionally, goal analysis consists of decomposing goals into subgoals through an AND- or OR-decomposition.
- Given a goal model and a set of initial labels for some goals (S for Datisfied and D for Denied) there is a simple labels propagation algorithm which can generate labels for all other goals of the model [Nilsson'72]

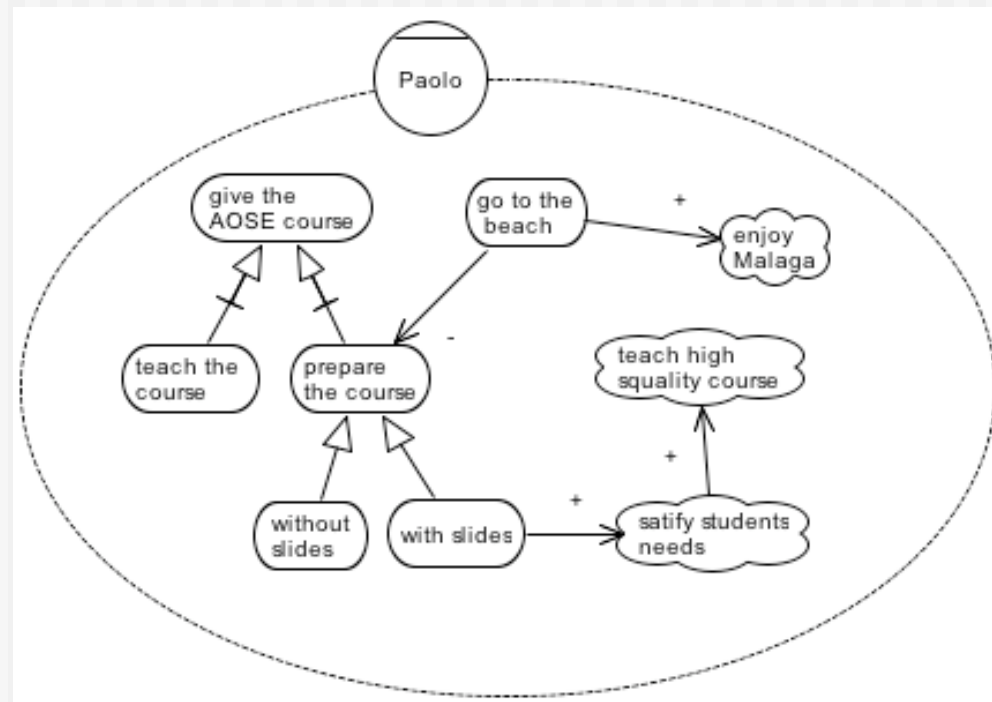


Goal Analysis

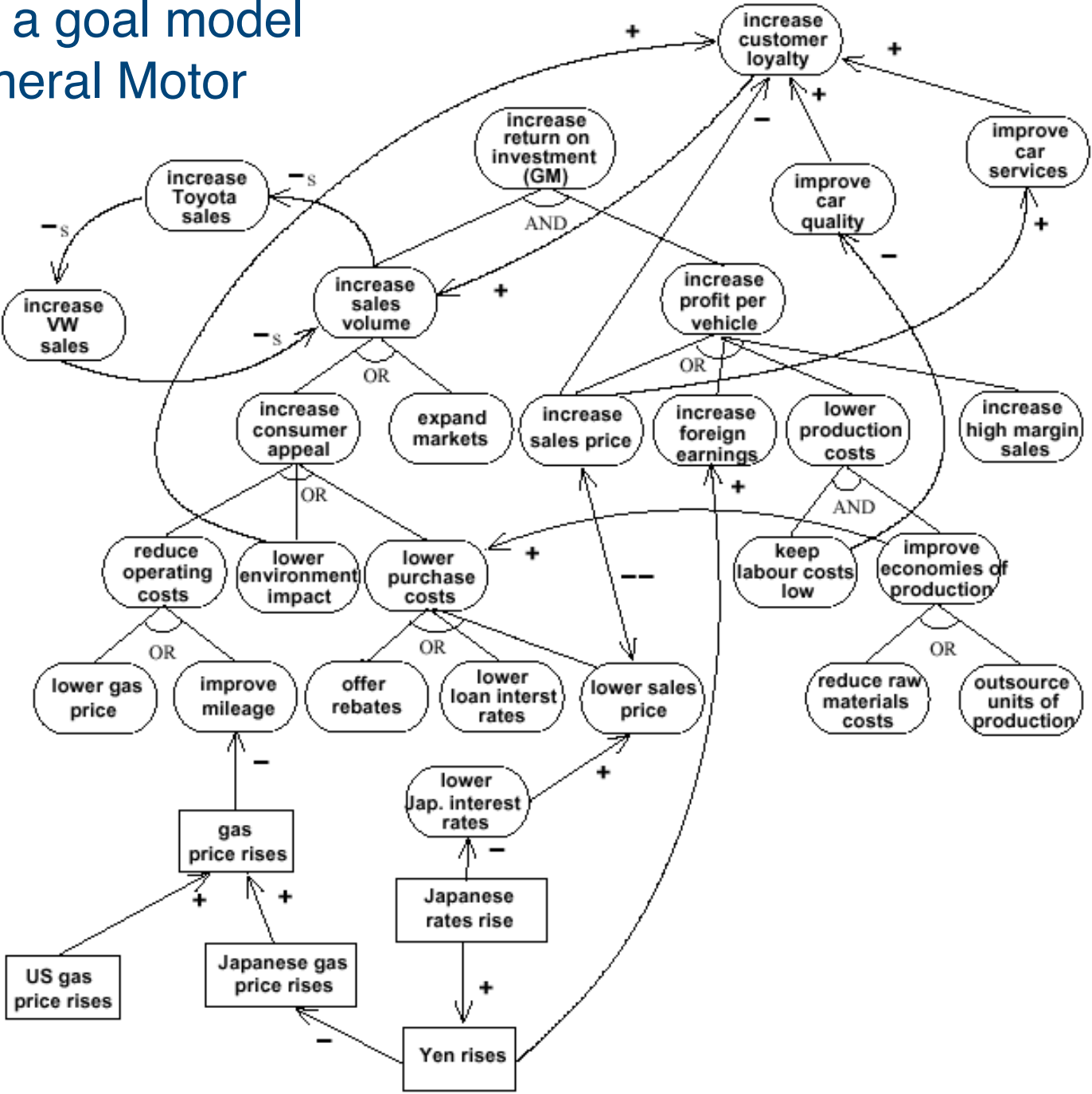
- However, there are many domains where goals are not formalizable and the relationships among them cannot be captured by semantically well-defined relations such as AND/OR ones.
- E.g., in RE:
 - `Highly reliable system`, has no formally defined predicate to prescribe its meaning, though you can define necessary conditions for its satisfaction
 - `Highly reliable system` can be related to other goals, such as `thoroughly debugged system` (the latter contributes to the satisfaction of the former) - partial and qualitative contribution

A simple example of goal model

- AND/OR decompositions
- Positive (+)/ Negative (-) contribution links



Part of a goal model for General Motor

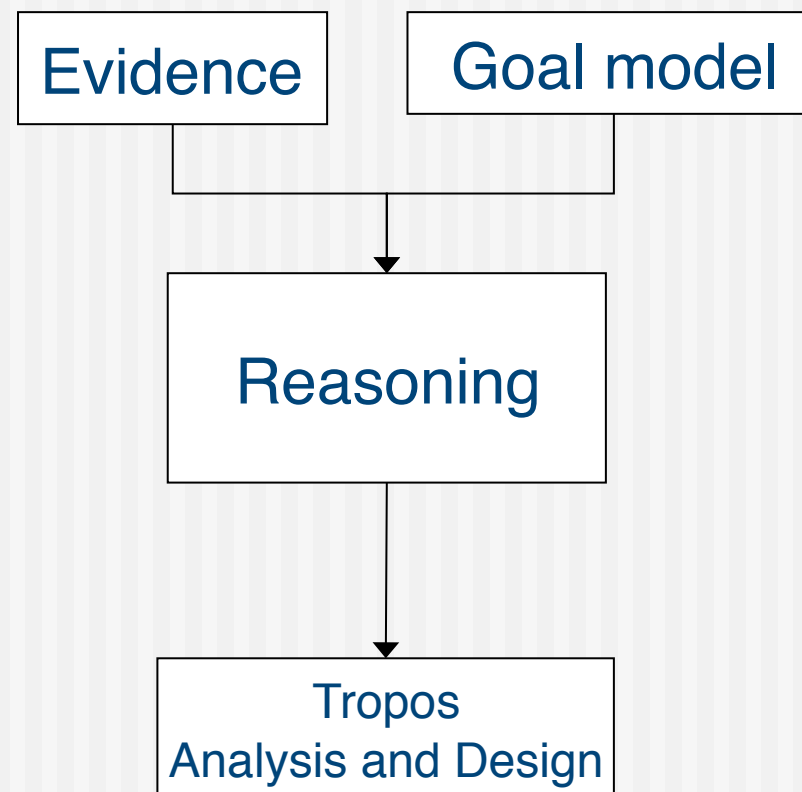


Goal modeling and reasoning

- In Tropos we need to capture relations among goals/softgoals
 - From early requirements analysis up to architectural design
- We need also to reason about satisfaction or denial of goals
 - If all subgoals are satisfied, top goals are satisfied
 - But what happen if a subgoal is partially satisfied?
 - ... and what happen if two goals are in conflict (E.g., `prepare the AOSE course` and `go to the beach`)?
- Different form of reasoning
 - What the minimal set of subgoals that allow me to satisfy all top goals?
 - If I satisfy a specific subset of leaf goals what happen to my top goals?
 - Qualitative reasoning (E.g., a goal is partially satisfied)
 - Quantitative reasoning (E.g, the probability for a goal to be satisfied)

The Tropos approach

- Evidence about satisfaction/denial of goals
- Reasoning mechanisms to propagate evidence in the model
- The reasoning output is used to support the analysis and design process



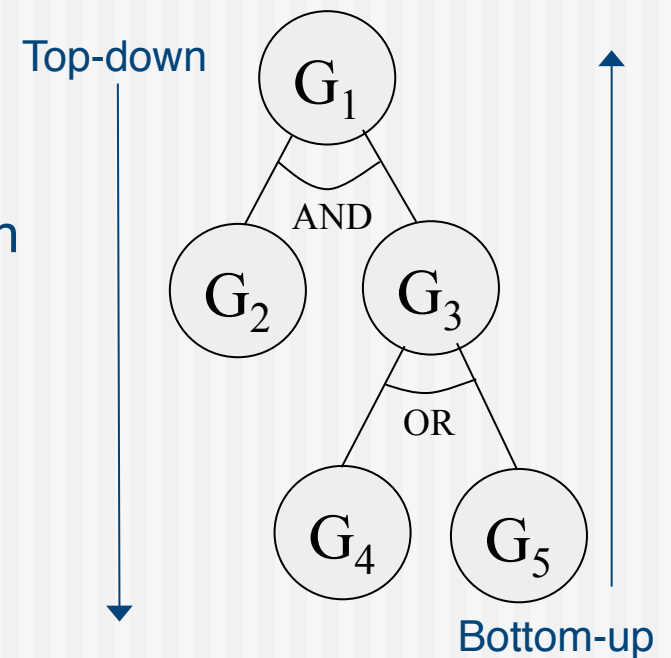
Goal Models

- Goal Dependency Graph:
 - Goals represented as Nodes
 - And/or relationships as (grouped) and/or arcs
 - Identity/negation as ++/- - arcs
 - Positive/negative contribution as +/- arcs
 - Cycles possible!
- Goal Valuations:
 - Goals can be either satisfied or denied
 - need to represent evidence of satisfaction/denial
 - Relationships propagate satisfaction and denial values
 - Conflicts possible!

The problem

Provide:

- Formal representation(s) of goal models and goal valuations
 - Qualitative and quantitative approach
- Formal techniques to reason on goal values and on their propagation through goal models
 - Top-down (backward) reasoning
 - Bottom-up (forward) reasoning



Qualitative approach

- Four predicates:
 - FS(g): there is at least Full evidence that g is Satisfied
 - PS(g): there is at least Partial evidence that g is Satisfied
 - FD(g): there is at least Full evidence that g is Denied
 - PD(g): there is at least Partial evidence that g is Denied
- Negated atoms \neg FS(g), \neg FD(g) not admitted!
- FS(g)/PS(g) independent from FD(g)/PD(g)
 - This allow to have conflicts
 - E.g., g can be fully satisfied and partially denied
 - Different sources of information can provide both evidence for satisfaction and denial
 - A goal can receive a negative contribution from another goal (you cannot do both!) but its actual decomposition allow to satisfy the goal

Axiomatization

- Axioms allow to capture (define) the semantics of goal models
 - Express the semantics of relations and value propagation
 - Used to build sound reasoning techniques

Goal	Invariant Axioms
g	$FS(g) \rightarrow PS(g)$ $FD(g) \rightarrow PD(g)$

Axiomatization

Goal Relation

Relation Axioms

$(G2, G3) \xrightarrow{AND} G1:$

$(FS(G2) \wedge FS(G3)) \rightarrow FS(G1)$
 $(PS(G2) \wedge PS(G3)) \rightarrow PS(G1)$
 $FD(G2) \rightarrow FD(G1), FD(G3) \rightarrow FD(G1)$
 $PD(G2) \rightarrow PD(G1), PD(G3) \rightarrow PD(G1)$

$(G2, G3) \xrightarrow{OR} G1:$

$(FS(G2) \vee FS(G3)) \rightarrow FS(G1)$
 $(PS(G2) \vee PS(G3)) \rightarrow PS(G1)$
 $FD(G2) \rightarrow FD(G1), FD(G3) \rightarrow FD(G1)$
 $PD(G2) \rightarrow PD(G1), PD(G3) \rightarrow PD(G1)$

Axiomatization

Goal Relation

Relation Axioms

$G_2 \xrightarrow{++S} G_1:$ $FS(G_2) \rightarrow FS(G_1), PS(G_2) \rightarrow PS(G_1)$

$G_2 \xrightarrow{--S} G_1:$ $FS(G_2) \rightarrow FD(G_1), PS(G_2) \rightarrow PD(G_1)$

$G_2 \xrightarrow{+S} G_1:$ $FS(G_2) \rightarrow PS(G_1), PS(G_2) \rightarrow PS(G_1)$

$G_2 \xrightarrow{-S} G_1:$ $FS(G_2) \rightarrow PD(G_1), PS(G_2) \rightarrow PD(G_1)$

$G_2 \xrightarrow{++D} G_1:$ $FD(G_2) \rightarrow FD(G_1), PD(G_2) \rightarrow PD(G_1)$

$G_2 \xrightarrow{--D} G_1:$ $FD(G_2) \rightarrow FS(G_1), PD(G_2) \rightarrow PS(G_1)$

$G_2 \xrightarrow{+D} G_1:$ $FD(G_2) \rightarrow PD(G_1), PD(G_2) \rightarrow PD(G_1)$

$G_2 \xrightarrow{-D} G_1:$ $FD(G_2) \rightarrow PS(G_1), PD(G_2) \rightarrow PS(G_1)$

Axiomatization (cont.)

- or, +D, -D, ++D, --D are dual w.r.t. and, +S, -S, ++S, --S
- Propagation of satisfaction through a ++, --, +, - may be or may be not symmetric w.r.t. that of denial:

$$\begin{aligned} G2 \xrightarrow{+} G1 &\Leftrightarrow G2 \xrightarrow{+S} G1 \text{ and } G2 \xrightarrow{+D} G1 \\ G2 \xrightarrow{-} G1 &\Leftrightarrow G2 \xrightarrow{-S} G1 \text{ and } G2 \xrightarrow{-D} G1 \end{aligned}$$

Satisfaction/Dianial:

g is totally satisfied [resp. partially satisfied, totally/partially denied] iff FS(g) [resp. PS(g), FD(g), PD(g)] can be logically inferred from the initial assignment and the axioms

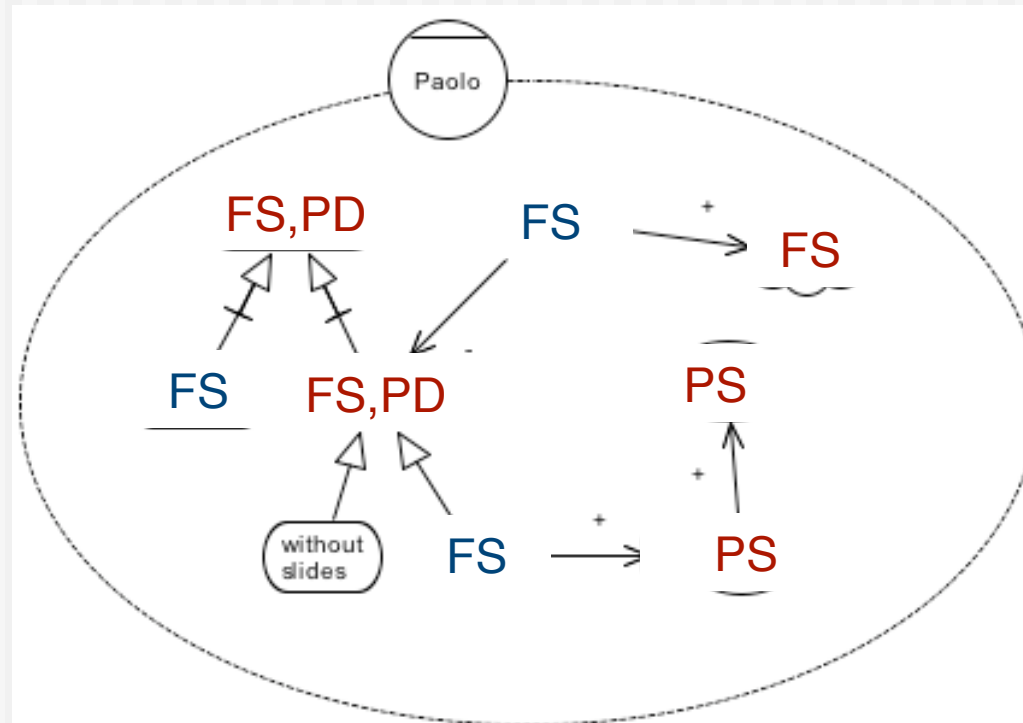
Forward Reasoning

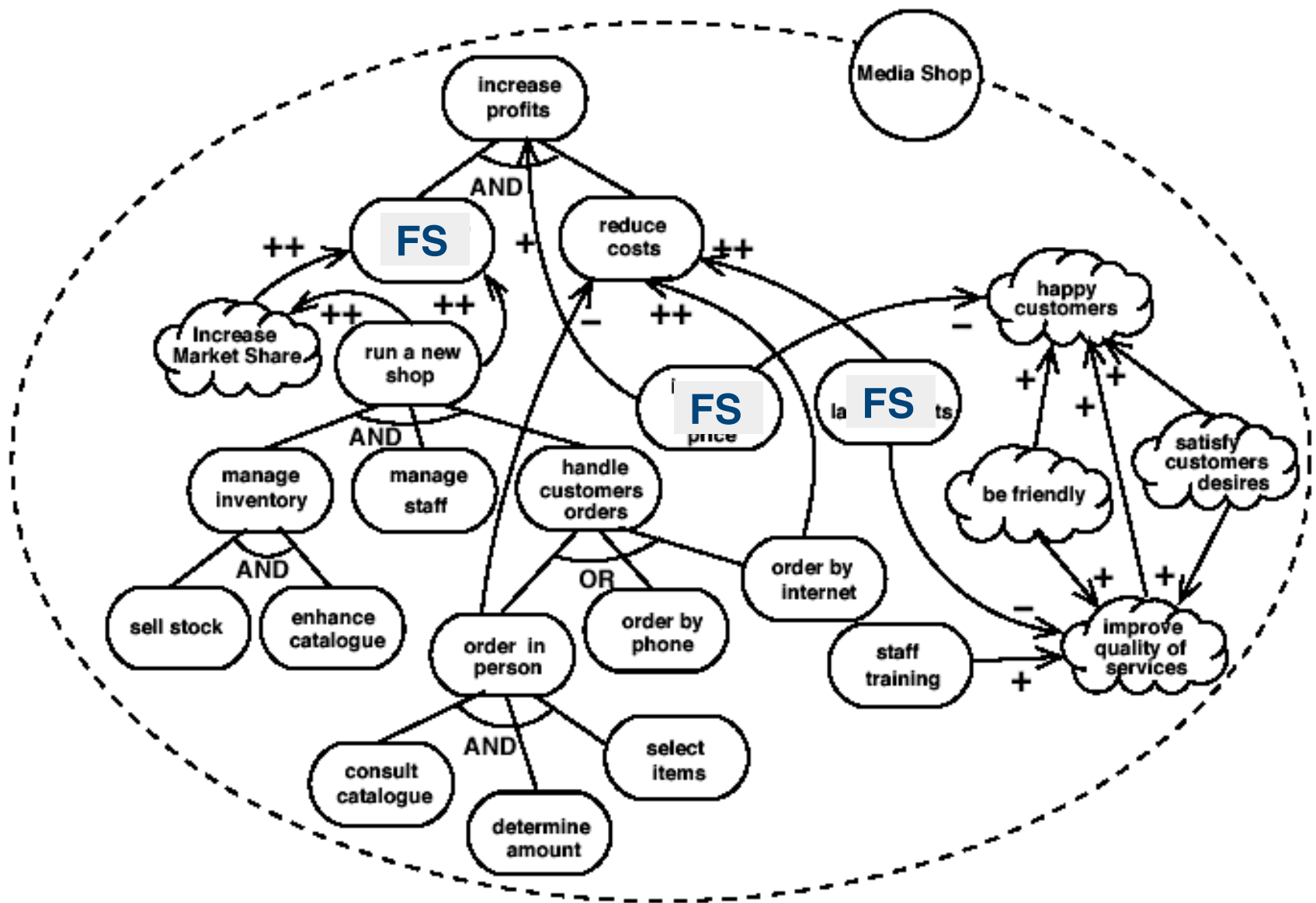
- Given
 - goal model
 - initial values assignment to some goals
(input goals -- typically leaf goals)
- Forward reasoning focuses on the forward propagation of these initial values to all other goals of the graph accordingly to the axioms
- Initial values represents the evidence (possibly contradictory) available about the satisfaction and the denial of goals: $\{FS(G1), PD(G2), \dots\}$
 - Usually provided by the domain expert(s)

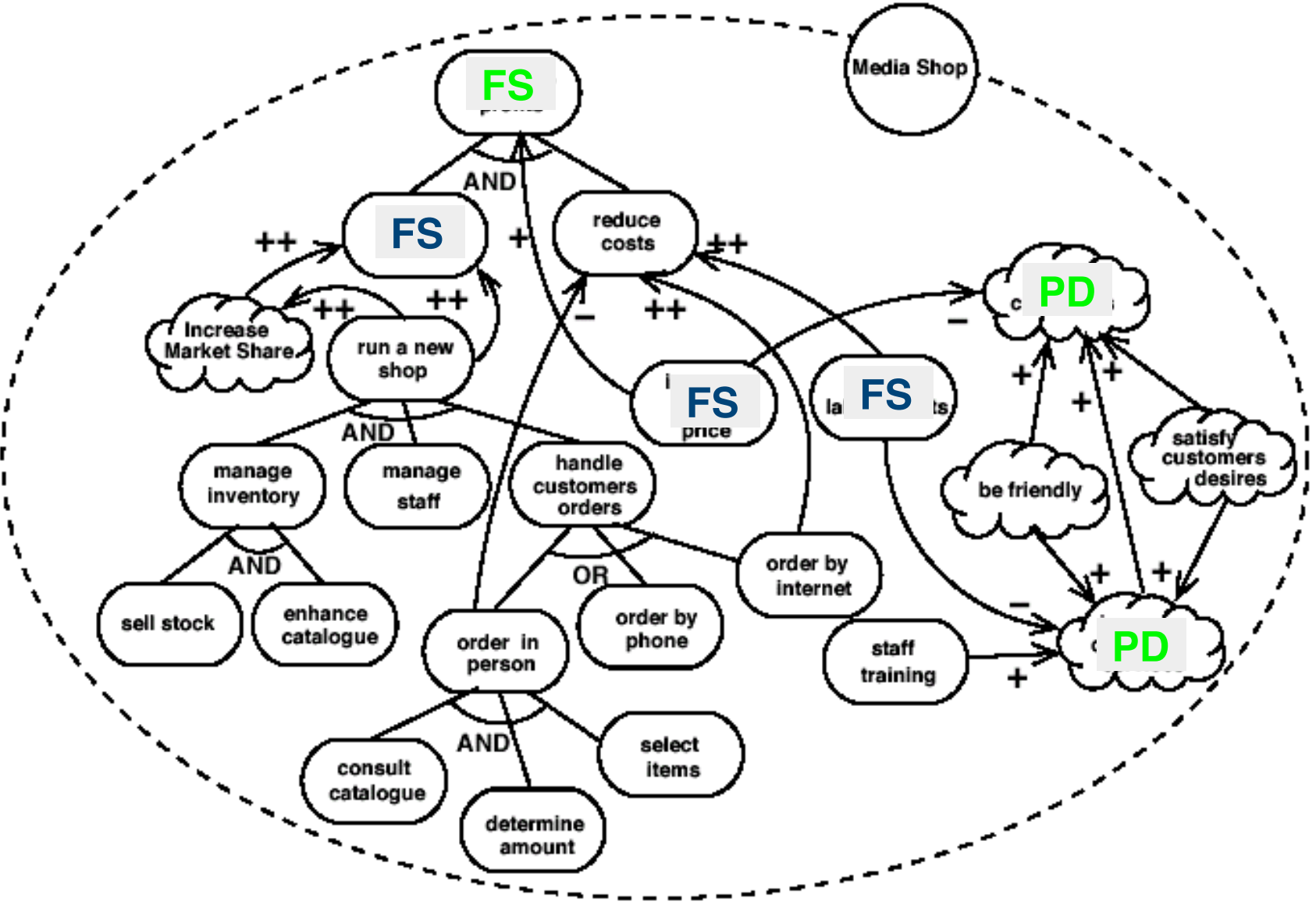
An example of Forward prop.

Initial: {FS(*teach the course*), FS(*with slides*), FS(*go to the beach*)}

Final : {FS(*teach the course*), FS(*with slides*), FS(*go to the beach*), FS(*prepare the course*), PD(*prepare the course*), PS(*enjoy Malaga*), PS(*satisfy students needs*), PS(*teach high quality course*).....}







Propagation Algorithm

```
1.  label_array Label_Graph(graph <G,R>,label_array Initail)
2.    Current=Initial;
3.    do
4.      Old=Current;
5.      for each  $G_i \in G$  do
6.        Current[i]=Update_label(i,<G,R>,Old);
7.    until not (Current==Old);
8.    return Current;

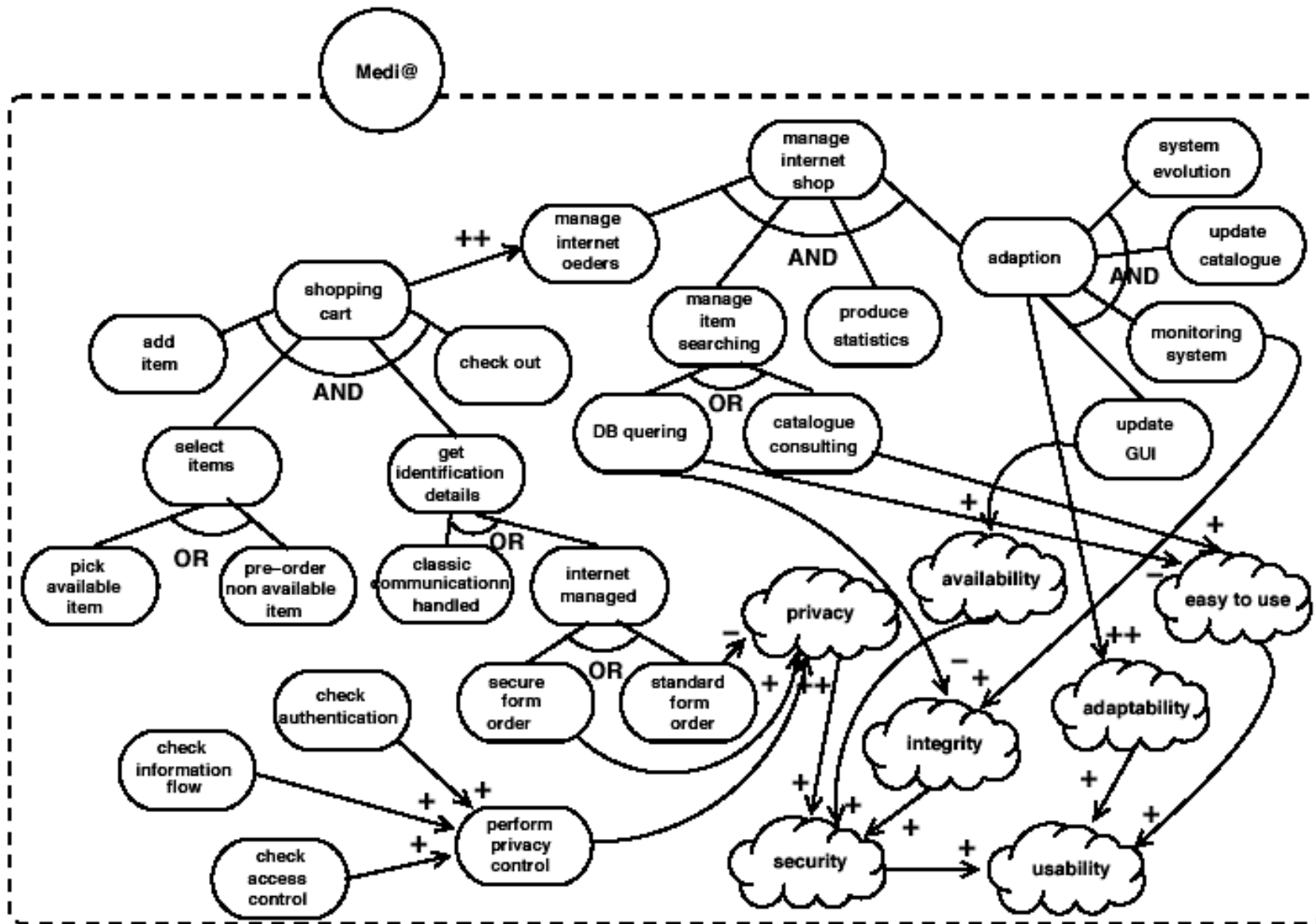
17. label Update_label(int i, graph <G,R>,label_array Old)
18.   for each  $R_j \in R$  s.t. target(Ri)==  $G_i$  do
19.     satij = Apply_Rules_Sat( $G_i, R_j, Old$ )
20.     denij = Apply_Rules_Den( $G_i, R_j, Old$ )
21.   return <max(maxj(satij), Old[i].sat),
           max(maxj(denij), Old[i].den)>
```

Propagation Algorithm (cont.)

- To each g we associate two variables $Sat(g)$ and $Den(g)$ ranging in $\{F,P,N\}$ such that $F > P > N$
- E.g., $Sat(g) \geq P$ states that there is at least partial evidence that g is satisfiable
- From the initial assignment, we propagate the values according to the following rules:

	$(G_2, G_3) \xrightarrow{and} G_1$	$G_2 \xrightarrow{+} G_1$	$G_2 \xrightarrow{-} G_1$	$G_2 \xrightarrow{++} G_1$	$G_2 \xrightarrow{--} G_1$
$Sat(G_1)$	$\min\{Sat(G_2), Sat(G_3)\}$	$\min\{Sat(G_2), P\}$	N	$Sat(G_2)$	N
$Den(G_1)$	$\max\{Sat(G_2), Sat(G_3)\}$	N	$\min\{Sat(G_2), P\}$	N	$Sat(G_2)$

- or, +D, -D, ++D, --D dual w.r.t. and, +S, -S, ++S, --S
- Satisfaction/denial values monotonically non-decreasing
- Terminates when reaches a fixpoint ($Current == Old$)



Goals	Exp 1		Exp 2		Exp 3		Exp 4	
	Init	Fin	Init	Fin	Init	Fin	Init	Fin
	S	D	S	D	S	D	S	D
DB querying	F		F					
catalogue consulting					F		F	
pick available item	F		F		F		F	
pre-order non available item								
classic communication handled	F							
standard form order			F		F			
secure form order							F	
manage internet shop		F		F		F		F
privacy		P		P P		P P		P
availability		P		P		P		P
integrity		P P		P		P P		P
usability		P		P		P		P
adaptability		F		F		F		F
easy to use		P		P		P		P
security		P		P		P		P

Forward Reasoning in Tropos

- It is adopted for evaluating the impact of the adoption of the different alternatives with respect to the
 - functional requirements (top goals)
 - non-functional requirements (softgoals)of the system-to-be
- Reasoning may involve
 - Single actor (intra actor reasoning)
 - Multiple actors (inter actor reasoning)

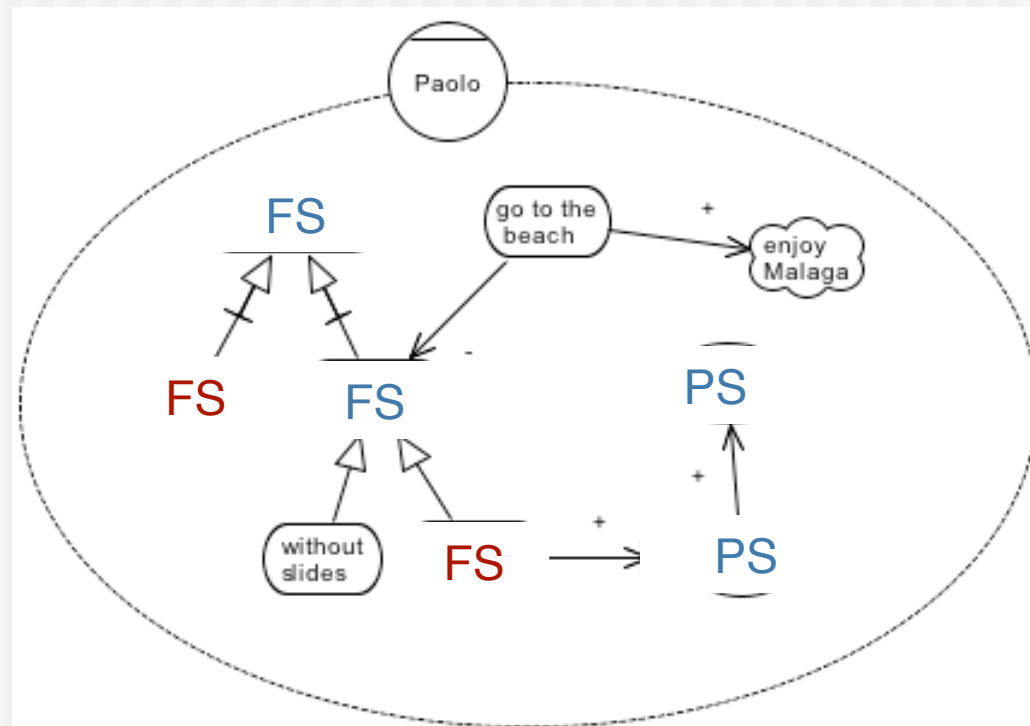
Backward reasoning

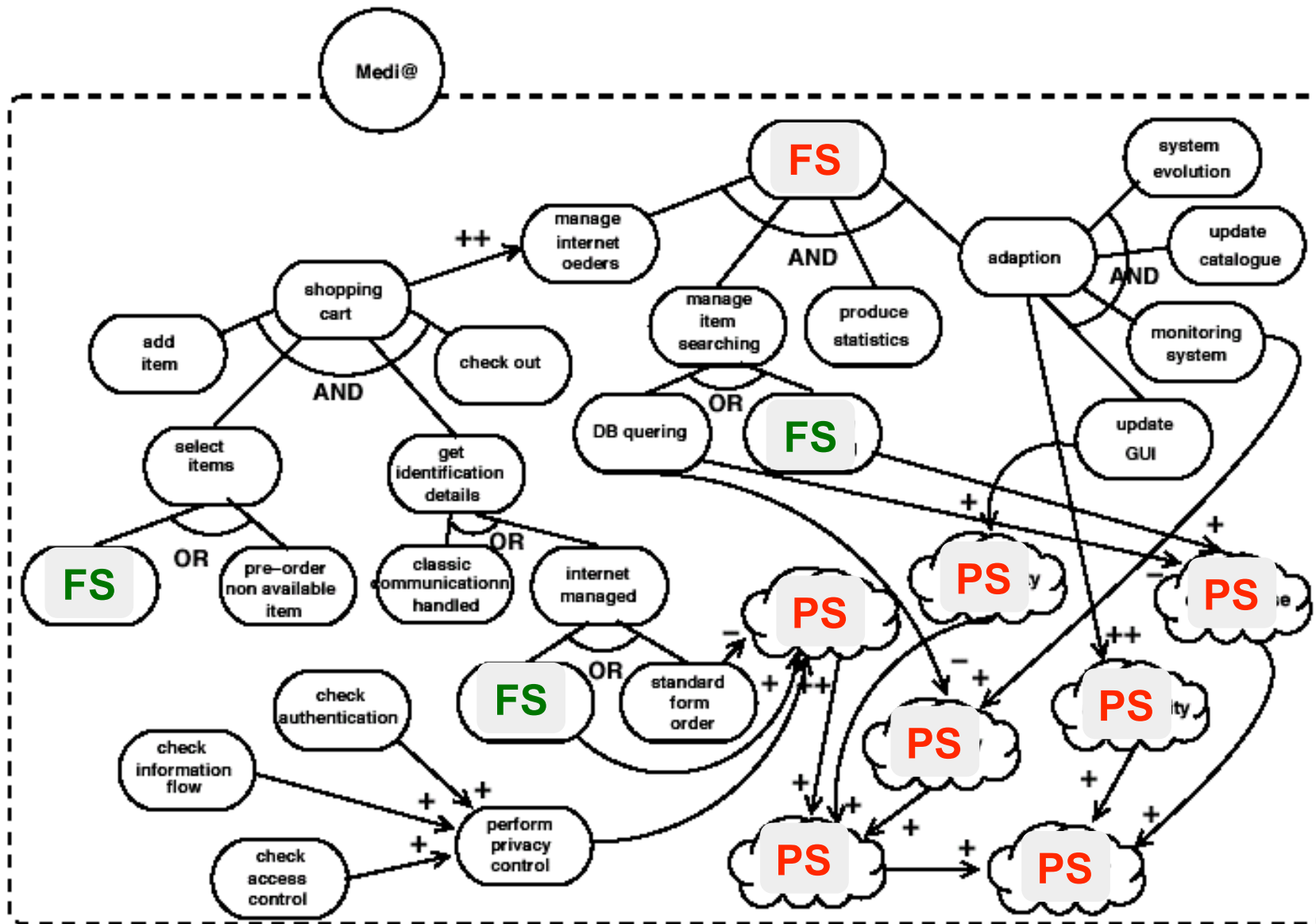
- We set the desired final values of the target goals, and we want to find possible initial assignments to the input goals which would cause the desired final values of the target goals
- We search for possible initial assignments to the input goals which would cause the desired final values of the target goals by forward propagation
- The user may also add some desired constraints, and decide to avoid conflicts

An example of Backward prop.

Final: {FS(*give AOSE course*), PS(*teach high quality course*)}

Assign. : {FS(*prepare the course*), PS(*with slides*)}





Constraints and Costs

- We may also add some desired constraints and decide to avoid
 - Strong conflict (e.g., FS(G),FD(G))
 - Medium conflict (e.g., FS(G),PD(G))
 - Weak (e.g., PS(G),PD(G))
 - all conflicts
- Assigning a cost to each input goal, we search for an assignment at the minimum cost

Propositional Satisfiability (SAT)

- We reduce the backward search to a SAT problem
- SAT is the problem of determining whether a boolean formula Φ admits at least one satisfying truth assignment μ to its variables A_i
- SAT is a NP-complete problem (there does not exist any polynomial algorithm able to solve it)
- There exists efficient SAT techniques
 - DPLL is the most popular SAT algorithm
 - **CHAFF** is the most efficient DPLL implementation
- There are several techniques to improve the efficiency of DPLL (e.g., backjumping, learning, random restart)

Minimum-Weight SAT (MW-SAT)

- MW-SAT is a variant of SAT, where the boolean variables A_i occurring in Φ are given a positive integer weight w_i
- MW-SAT is the problem of determining a truth assignment μ satisfying Φ which minimizes the value

$$W(\mu) := \sum_i \{w_i \mid A_i \text{ is assigned } \top \text{ by } \mu\}.$$

or stating there is none.

- **MINWEIGHT** is the state-of-art solver for MW-SAT

Basic Formalization

- The boolean variables of Φ are all the values $FS(G), PS(G), FD(G), PD(G)$ for each goal G and Φ is

$$\Phi := \Phi_{graph} \wedge \Phi_{outval} \wedge \Phi_{backward} [\wedge \Phi_{constraints} \wedge \Phi_{conflict}]$$

where

Φ_{graph} encodes the *goal graph*

Φ_{outval} encodes the *desired final output values*

$\Phi_{backward}$ encodes the *backward reasoning*

$\Phi_{constraints}$ encodes *user's constraints* (optional)

$\Phi_{conflict}$ encodes the *prevention of conflicts* (optional)

Basic Formalization cont.

- **Encoding the goal graph**

$$\Phi_{graph} := \bigwedge_{G \in \mathcal{G}} Invar_Ax(G) \wedge \bigwedge_{r \in \mathcal{R}} Rel_Ax(r).$$

$Invar_Ax(G)$ is the conjunction of the invariant axioms and $Rel_Ax(r)$ is the conjunction of the relation axioms (forward propagation through the relation arcs in the graph)

- **Representing Desired Final Output Values**

$$\Phi_{outval} := \bigwedge_{G \in Target(\mathcal{G})} vs(G) \wedge \bigwedge_{G \in Target(\mathcal{G})} vd(G)$$

$Target(\mathcal{G})$ is the set of target goals and $vs(G) \in \{T, PS(G), FS(G)\}$, $vd(G) \in \{T, PD(G), FD(G)\}$ are the maximum satisfiability and deniability values assigned to the target goal G .

Basic Formalization cont.

■ Encoding Backward Reasoning

$$\Phi_{backward} := \bigwedge_{G \in \mathcal{G}/Input(\mathcal{G})} \bigwedge_{v(G)} Backward_Ax(v(G))$$
$$Backward_Ax(v(G)) := v(G) \rightarrow \bigvee_{r \in Incoming(G)} Prereqs(v(G), r)$$

$Input(\mathcal{G})$ is the set of input goals; $Incoming(G)$ is the set of relations in \mathcal{G} ; $v(G) = \{PS(G), FS(G), PD(G), FD(G)\}$, and $Prereqs(v(G), r)$ is a formula which is true iff the prerequisites of $v(G)$ through r hold.

$Backward_Ax(v(G))$ is the set of propagation axioms (see next slide)

If G is not an input goal and $v(G)$ holds, then this value must derive from the prerequisite values of some incoming relations of \mathcal{G}

Axioms for backward propagation

$$\begin{array}{l}
 PS(G) \rightarrow \left(\begin{array}{ll}
 \bigwedge_i PS(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{and}} G \\
 \bigvee_i PS(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{or}} G \\
 PS(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{+S} G \\
 FD(G_i) & \text{For every } R_i: G_i \xrightarrow{-D} G
 \end{array} \right. \\
 \\
 PS(G) \rightarrow \left(\begin{array}{ll}
 \bigwedge_i PS(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{and}} G \\
 \bigvee_i PS(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{or}} G \\
 PS(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{+S} G \\
 PD(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{-D} G \\
 PS(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{+S} G \\
 PD(G_i) & \text{For every } R_i: G_i \xrightarrow{-D} G
 \end{array} \right. \\
 \\
 FD(G) \rightarrow \left(\begin{array}{ll}
 \bigvee_i FD(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{and}} G \\
 \bigwedge_i FD(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{or}} G \\
 FD(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{+D} G \\
 PS(G_i) & \text{For every } R_i: G_i \xrightarrow{-S} G
 \end{array} \right. \\
 \\
 PD(G) \rightarrow \left(\begin{array}{ll}
 \bigvee_i PD(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{and}} G \\
 \bigwedge_i PD(G_i) \vee & \text{If } (G_1, \dots, G_i, \dots, G_n) \xrightarrow{\text{or}} G \\
 PD(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{+D} G \\
 PS(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{-S} G \\
 PD(G_i) \vee & \text{For every } R_i: G_i \xrightarrow{+D} G \\
 PS(G_i) & \text{For every } R_i: G_i \xrightarrow{-S} G
 \end{array} \right.
 \end{array}$$

Optional components

Adding User's Constraints and Desiderata

$$\Phi_{outval} := \bigwedge_i \bigvee_j lit_{ij}$$

The user expresses constraints and desiderata on goal values (e.g., “PS(G1)” means “G1 is at least partial satisfiable”, but it might totally satisfiable)

A negative clause value is used to prevent a value to a goal (e.g., “¬FD(G1)” means “G1 cannot be fully deniable”, but it might be partially deniable)

FS(G1)∨FS(G2) means at least G1 or G2 must be fully satisfiable

Basic Formalization cont.

Preventing conflicts

It allows the user for looking for solutions which do not involve conflicts

Strong conflicts

$$\Phi_{conflict} := \bigwedge_{G \in \mathcal{G}} (\neg FS(G) \vee \neg FD(G))$$

Strong and medium conflicts

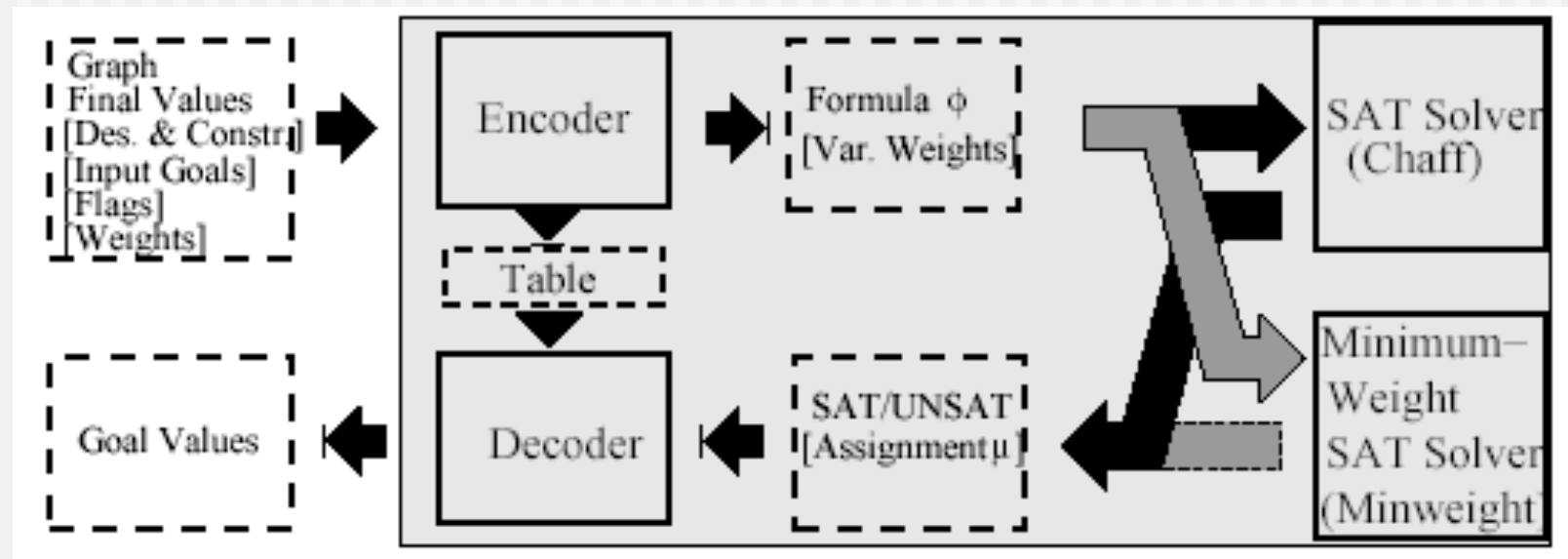
$$\Phi_{conflict} := \bigwedge_{G \in \mathcal{G}} ((\neg FS(G) \vee \neg PD(G)) \wedge (\neg PS(G) \vee \neg FD(G)))$$

All conflicts

$$\Phi_{conflict} := \bigwedge_{G \in \mathcal{G}} (\neg PS(G) \vee \neg PD(G))$$

Backward Prop. implementation

- We have reduced the qualitative problem to the Satisfiability (SAT) and minimum-cost satisfiability (minimum-cost SAT) problems for Boolean formulas
- GOLSOLVE / GOLMINSOLVE



Backward reasoning in Tropos

- Used to find the set of goals at the minimum costs that if achieved they can guarantee the achievement of the desired top goals (functional requirements) and softgoals (non-functional requirements).
- In other words, we find among the alternatives of the goal model those with the minimal cost that allow us to obtain our desired goals.
- Reasoning may involve
 - Single actor (intra actor reasoning)
 - Multiple actors (inter actor reasoning)

C=0

C=16

C=14

C=13

Goals	Exp 1		Exp 2		Exp 3		Exp 4	
	Init	Fin	Init	Fin	Init	Fin	Init	Fin
	S	D	S	D	S	D	S	D
DB querying (3)								F
catalogue consulting (6)				F		F		
pick available item (2)				F		F		F
pre-order non available item (7)								
classic communication handled (4)								
standard form order (6)						F		
secure form order (8)				F				F
manage internet shop	F		F	F	F	F	F	F
privacy	F		P	P	P	P	P	P
availability	F		P	P	P	P		P
integrity	F		P	P	P	P		P
usability	F		P	P	P	P		P
adaptability	F		P	F	P	F		F
easy to use	F		P	P	P	P		P
security	F		P	P	P	P		P

Avoiding conflicts

Quantitative approach

Quantitative Approach

- Evidence of satisfaction/denial represented by real values in \mathcal{D} :
 $[inf, sup]$, $0 \leq inf < sup$
- Value propagation through goal graphs as math functions, $f : \mathcal{D}^n \rightarrow \mathcal{D}$
- Much finer-grained:
 - Different degrees of satisfaction/denial evidence
 - Different degrees of positive/negative contribution
 - Different strength of conflicts

Numerical representation of evidence

- $Sat(g), Den(g) \in [inf, sup]$
- Atoms in the form $Sat(g) \geq c1$ [$Den(g) \geq c2$]: “there is at least an evidence $c1$ [$c2$] that g is Satisfied [Denied]”

$$c1 = inf, c2 = inf \Leftrightarrow \top$$

$$c1, c2 \in]inf, sup[\Leftrightarrow PS(g), PD(g)$$

$$c1 = sup, c2 = sup \Leftrightarrow FS(g), FD(g)$$

- *Conflict*: $Sat(g) \geq c1$ and $Den(g) \geq c2$, $c1, c2 \in]inf, sup]$

Value propagation model

- 2 dual OPERATORS: \oplus and \otimes , representing value propagation through “or” and “and”

- Independent probability model:

$$inf=0, sup=1$$

$$p1 \oplus p2 = p1 + p2 - p1 \cdot p2$$

$$p1 \otimes p2 = p1 \cdot p2$$

(disjunction and conjunction of two independent events of probability $p1$ and $p2$)

- Flow model (Resistor):

$$inf=0, sup = +\infty$$

$$v1 \oplus v2 = v1 + v2$$

$$v1 \otimes v2 = (v1 \cdot v2) / (v1 + v2)$$

- ...

Axiomatization

Goal Relation

Relation Axioms

$(G2, G3) \xrightarrow{AND} G1:$	$(Sat(G2) \geq x \wedge Sat(G3) \geq y) \rightarrow Sat(G1) \geq (x \otimes y)$ $(Den(G2) \geq x \wedge Den(G3) \geq y) \rightarrow Den(G1) \geq (x \oplus y)$
$(G2, G3) \xrightarrow{OR} G1:$	$(Sat(G2) \geq x \vee Sat(G3) \geq y) \rightarrow Sat(G1) \geq (x \oplus y)$ $(Den(G2) \geq x \vee Den(G3) \geq y) \rightarrow Den(G1) \geq (x \otimes y)$

- AND and OR relation are dual

Axiomatization cont.

Goal Relation	Relation Axioms
$G2 \xrightarrow{w+S} G1:$	$Sat(G2) \geq x \rightarrow Sat(G1) \geq (x \otimes w)$
$G2 \xrightarrow{w-S} G1:$	$Sat(G2) \geq x \rightarrow Den(G1) \geq (x \otimes w)$
$G2 \xrightarrow{++S} G1:$	$Sat(G2) \geq x \rightarrow Sat(G1) \geq x$
$G2 \xrightarrow{--S} G1:$	$Sat(G2) \geq x \rightarrow Den(G1) \geq x$
$G2 \xrightarrow{w+D} G1:$	$Den(G2) \geq x \rightarrow Den(G1) \geq (x \otimes w)$
$G2 \xrightarrow{w-D} G1:$	$Den(G2) \geq x \rightarrow Sat(G1) \geq (x \otimes w)$
$G2 \xrightarrow{++D} G1:$	$Den(G2) \geq x \rightarrow Den(G1) \geq x$
$G2 \xrightarrow{--D} G1:$	$Den(G2) \geq x \rightarrow sat(G1) \geq x$

- +D, -D, ++D, --D dual w.r.t. +S, -S, ++S, --S
- Remark: + and - relations have a weight w

Quantitative propagation

- There is at least an evidence c that g is satisfied [resp. denied] iff $Sat(g) \geq c$ [resp. $Den(g) \geq c$] can be logically inferred from the initial assignment and the axioms.
- $Sat(g) \geq c, Den(g) \geq c$ propagated independently

Forward Propagation Algorithm

```
1. label_array Label_Graph(graph  $\langle \mathbf{G}, \mathbf{R} \rangle$ , label_array Initial)
2.   Current = Initial;
3.   do
4.     Old = Current;
5.     for each  $G_i \in \mathbf{G}$  do
6.       Current[i] = Update_label(i,  $\langle \mathbf{G}, \mathbf{R} \rangle$ , Old);
7.   until not ( $\| \mathbf{Current} - \mathbf{Old} \|_{\infty} \leq \epsilon$ );
8.   return Current;
```

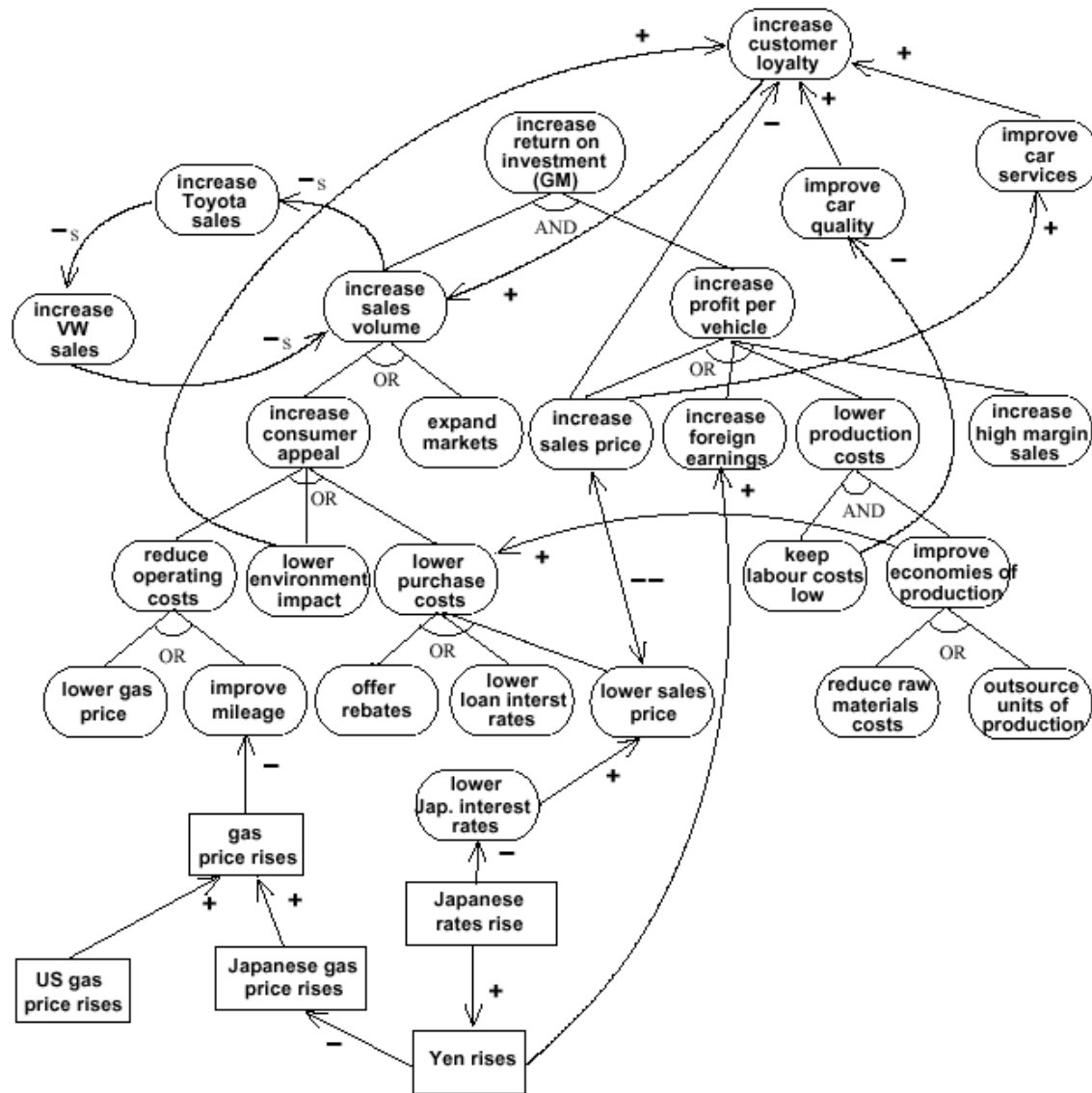
```
17. label Update_label(int i, graph  $\langle \mathbf{G}, \mathbf{R} \rangle$ , label_array Old)
18.   for each  $R_j \in \mathbf{R}$  s.t. target( $R_j$ ) ==  $G_i$  do
19.     satij = Apply_Rules_Sat( $G_i, R_j, Old$ )
20.     denij = Apply_Rules_Den( $G_i, R_j, Old$ )
21.   return  $\langle \max(\max_j(\text{sat}_{ij}), Old[i].\text{sat}),$   
            $\max(\max_j(\text{den}_{ij}), Old[i].\text{den}) \rangle$ 
```


Forward Propagation Algorithm

	$(G_2, G_3) \xrightarrow{and} G_1$	$G_2 \xrightarrow{w+} G_1$	$G_2 \xrightarrow{w-} G_1$	$G_2 \xrightarrow{++} G_1$	$G_2 \xrightarrow{--} G_1$
$Sat(G_1)$	$Sat(G_2) \otimes Sat(G_3)$	$Sat(G_2) \otimes w$		$Sat(G_2)$	
$Den(G_1)$	$Den(G_2) \oplus Den(G_3)$		$Sat(G_2) \otimes w$		$Sat(G_2)$

- Satisfaction/denial values monotonically non-decreasing
- Uses Cauchy-convergence as termination condition:

$$|a_{n+1} - a_n| \xrightarrow{n \rightarrow \infty} 0$$



Quantitative approach: example

Goal/Event	Relationship	Goal/Event
increase sales volume	0.6_{-S}	increase Toyota sales
increase Toyota sales	0.6_{-S}	increase VW sales
increase VW sales	0.6_{-S}	increase sales volume
increase customer loyalty	0.4_{+}	increase sales volume
increase sales prices	0.5_{-}	increase customer loyalty
increase car quality	0.8_{+}	increase customer loyalty
improve car services	0.7_{+}	increase customer loyalty
lower environment impact	0.4_{+}	increase customer loyalty
increase sales prices	0.3_{+}	improve car services
keep labour costs low	0.7_{-}	increase car quality
improve economies of production	0.8_{+}	lower purchase costs
Yen rises	0.8_{+}	increase foreign earnings
lower Japanese interest rates	0.4_{+}	lower sales price
Japanese rates rises	0.8_{-}	lower Japanese interest rates
Japanese rates rises	0.6_{+}	Yen rises
Yen rises	0.4_{-}	Japanese gas price rises
Japanese gas price rises	0.6_{+}	gas price rises
US gas price rises	0.6_{+}	gas price rises
gas price rises	0.8_{-}	improve mileage

Backward Propagation

- Formalization of the problem:
 - Linear cost function: $\min Ax$
 - a set of non linear equality and inequality constraints
- ... and we pass the system to Lingo 8.0.
- www.lindo.com

G(non-input):

$$Sat(G) = \max \begin{cases} \bigotimes_{i=1}^n Sat(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{and} G \\ \bigoplus_{i=1}^n Sat(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{or} G \\ \max(Sat(G_i) \cdot w), & \text{per ogni relazione } R_i : G_i \mapsto^{+S} G \\ \max(Sat(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{++S} G \\ \max(Den(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{--D} G \\ \max(Den(G_i) \cdot w), & \text{per ogni relazione } R : G_i \mapsto^{-D} G \end{cases}$$

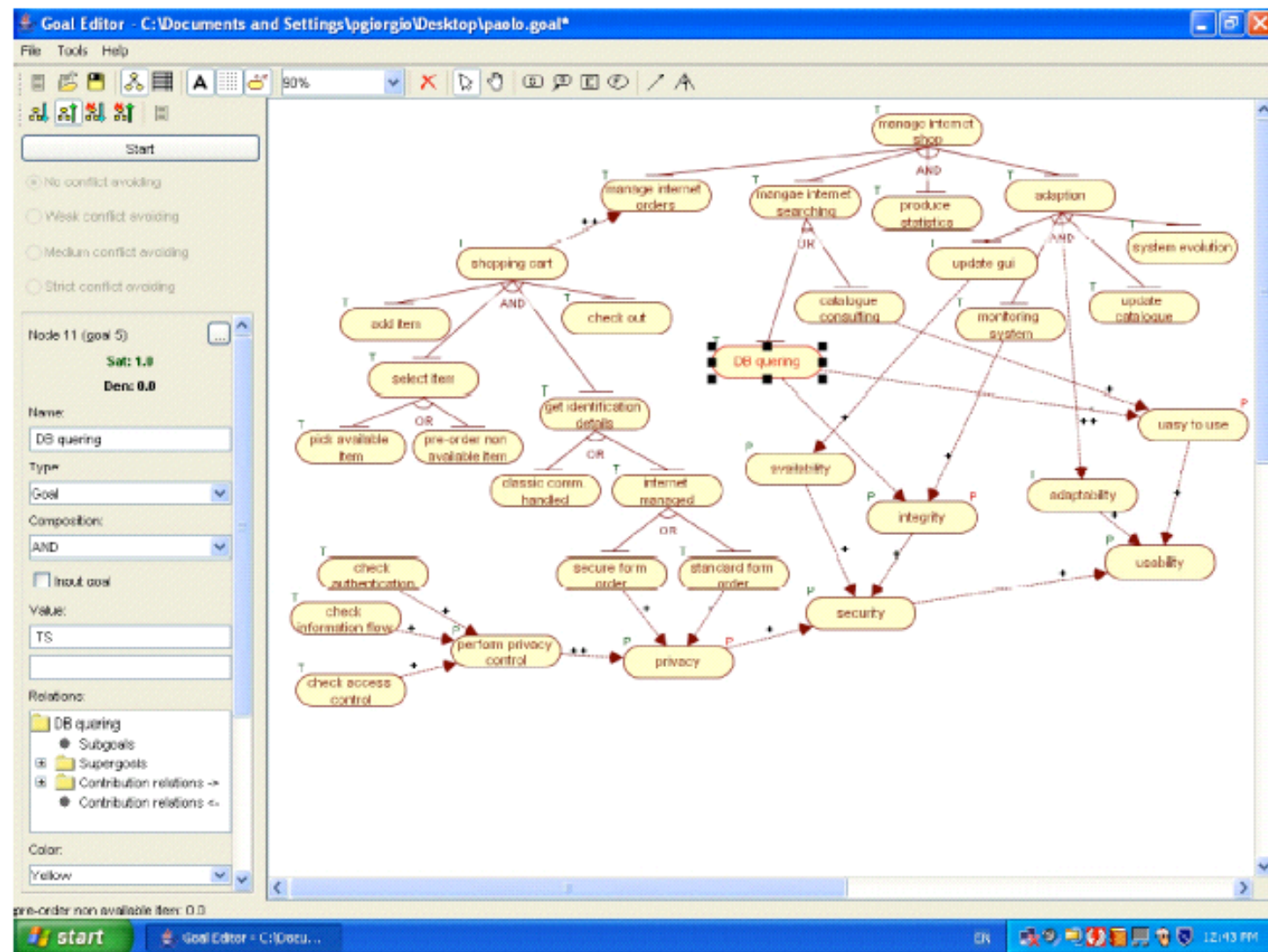
$$Den(G) = \max \begin{cases} \bigoplus_{i=1}^n Den(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{and} G \\ \bigotimes_{i=1}^n Den(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{or} G \\ \max(Den(G_i) \cdot w), & \text{per ogni relazione } R_i : G_i \mapsto^{+D} G \\ \max(Den(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{++D} G \\ \max(Sat(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{--S} G \\ \max(Sat(G_i) \cdot w), & \text{per ogni relazione } R : G_i \mapsto^{-S} G \end{cases}$$

G(input):

$$Sat(G) \geq \max \begin{cases} \bigotimes_{i=1}^n Sat(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{and} G \\ \bigoplus_{i=1}^n Sat(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{or} G \\ \max(Sat(G_i) \cdot w), & \text{per ogni relazione } R_i : G_i \mapsto^{+S} G \\ \max(Sat(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{++S} G \\ \max(Den(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{--D} G \\ \max(Den(G_i) \cdot w), & \text{per ogni relazione } R : G_i \mapsto^{-D} G \end{cases}$$

$$Den(G) \geq \max \begin{cases} \bigoplus_{i=1}^n Den(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{and} G \\ \bigotimes_{i=1}^n Den(G_i), & \text{se } (G_1 \dots G_i \dots G_n) \mapsto^{or} G \\ \max(Den(G_i) \cdot w), & \text{per ogni relazione } R_i : G_i \mapsto^{+D} G \\ \max(Den(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{++D} G \\ \max(Sat(G_i)), & \text{per ogni relazione } R : G_i \mapsto^{--S} G \\ \max(Sat(G_i) \cdot w), & \text{per ogni relazione } R : G_i \mapsto^{-S} G \end{cases}$$

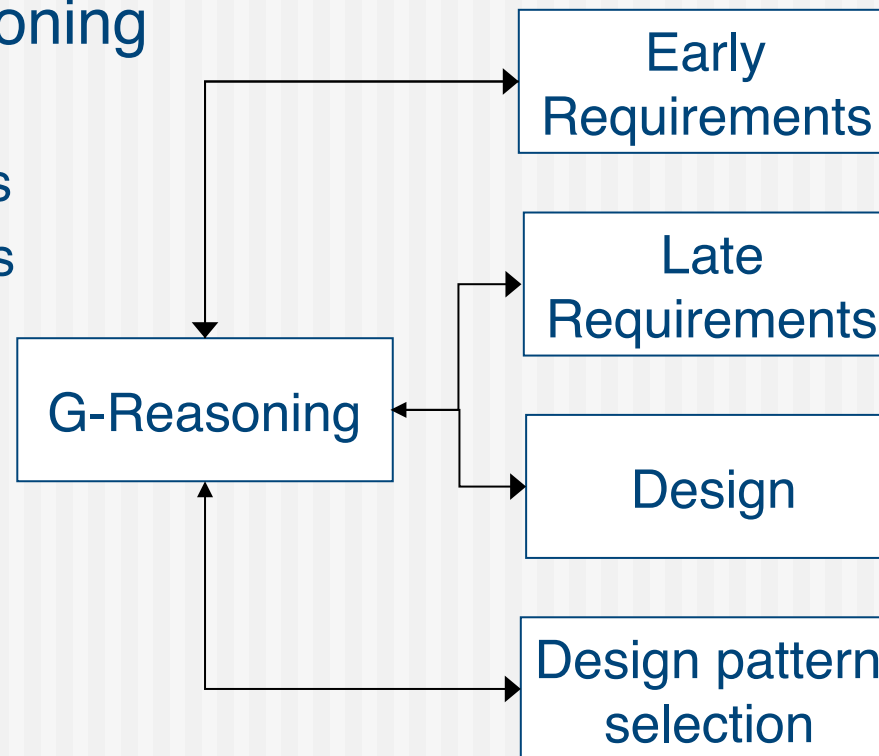
GR Tool



Reasoning about goal models

- Supporting reasoning

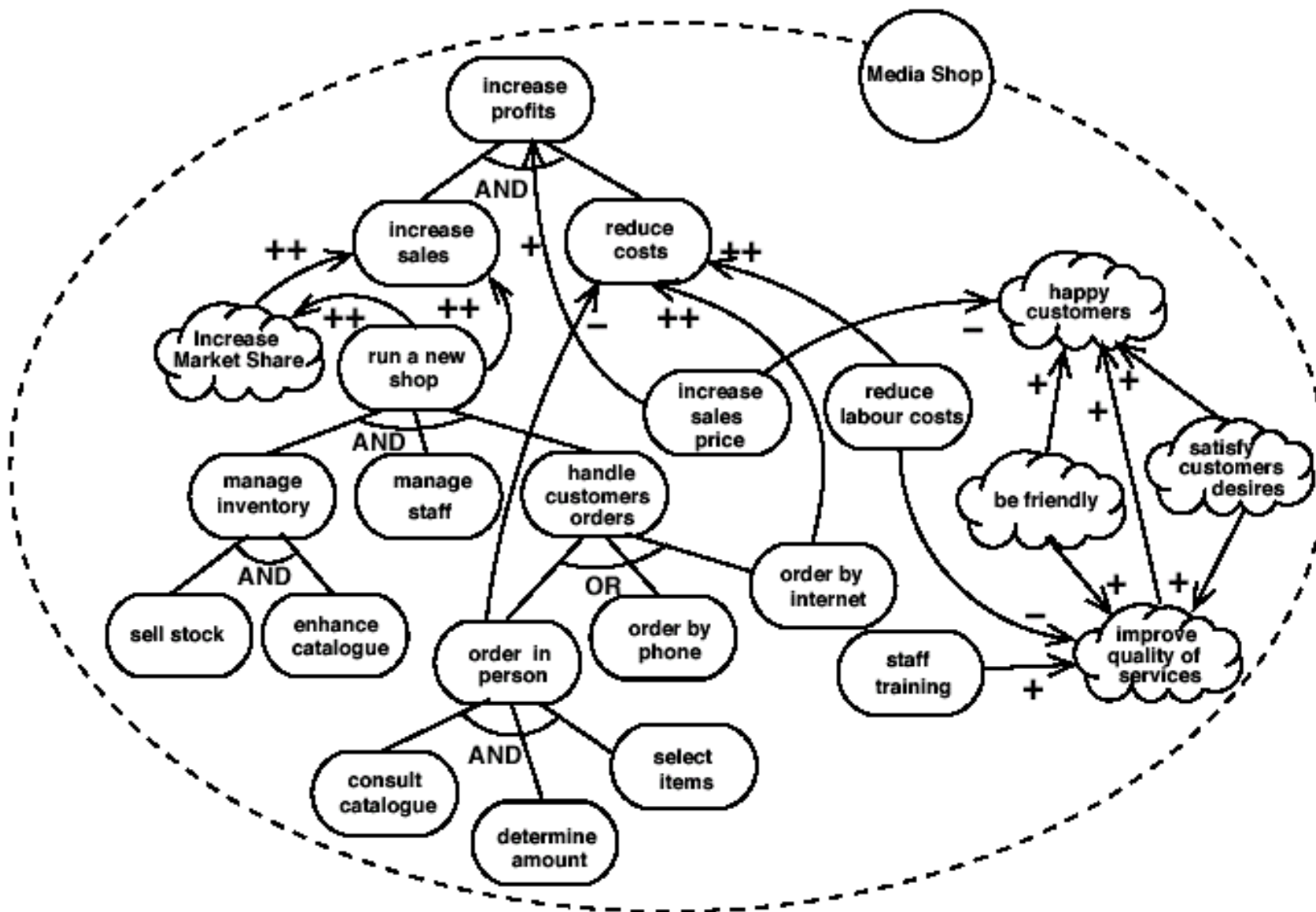
- Different uses
- Different models
- Different outputs



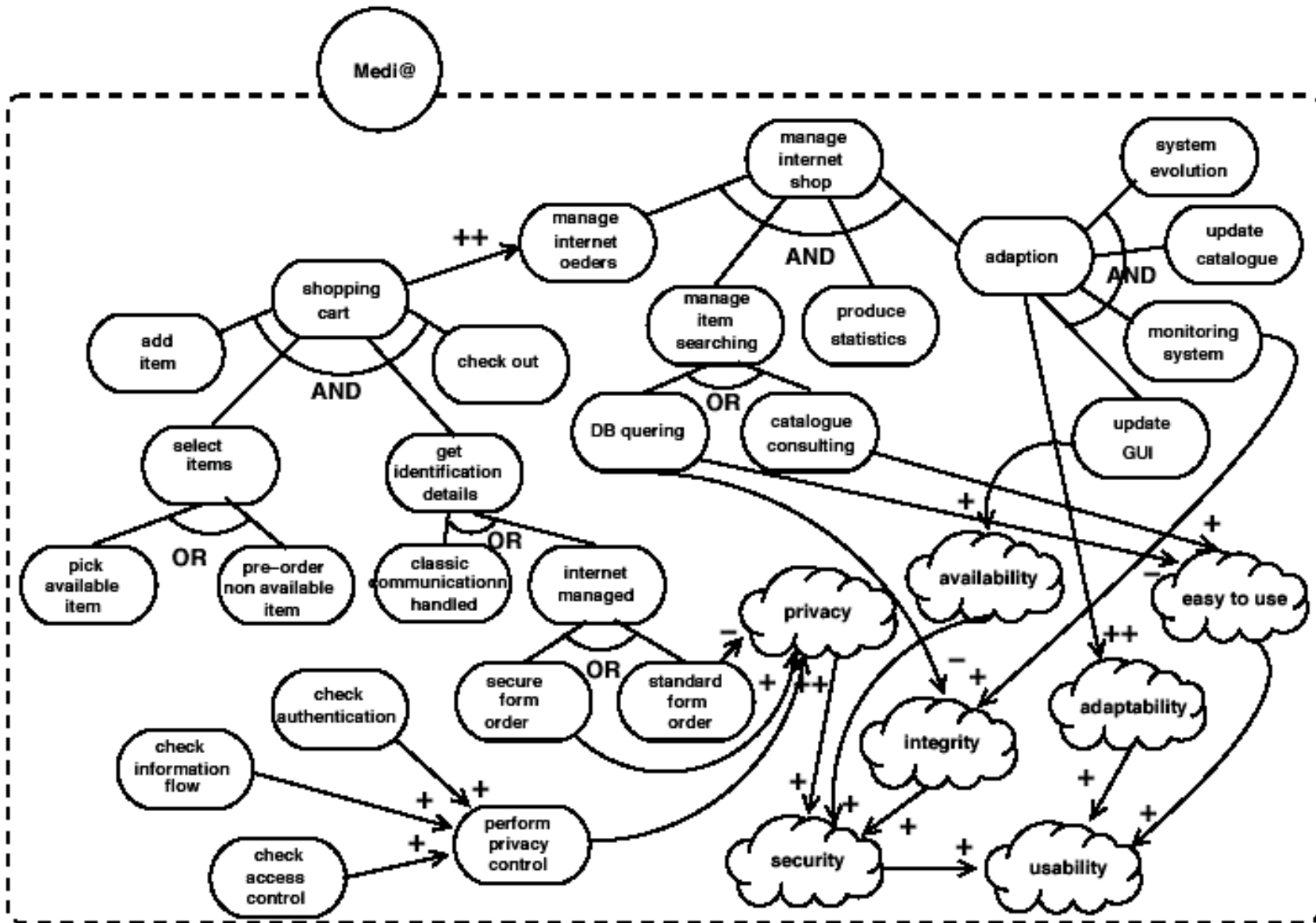
Reasoning about goal models cont.

- Early requirements analysis
 - Allow to analyze the organizational setting intra e inter actor analysis
 - Verify satisfaction of goals
 - A means to discuss with the stakeholders about their goals
 - Find possible conflicts
- Late requirements analysis
 - Evaluate possible alternative functional requirements wrt non-functional requirements (softgoals)
 - Find possible conflicts among requirements
 - Reason about requirements impact over stakeholders goals/softgoals

Actor diagram: *Early requirements analysis*

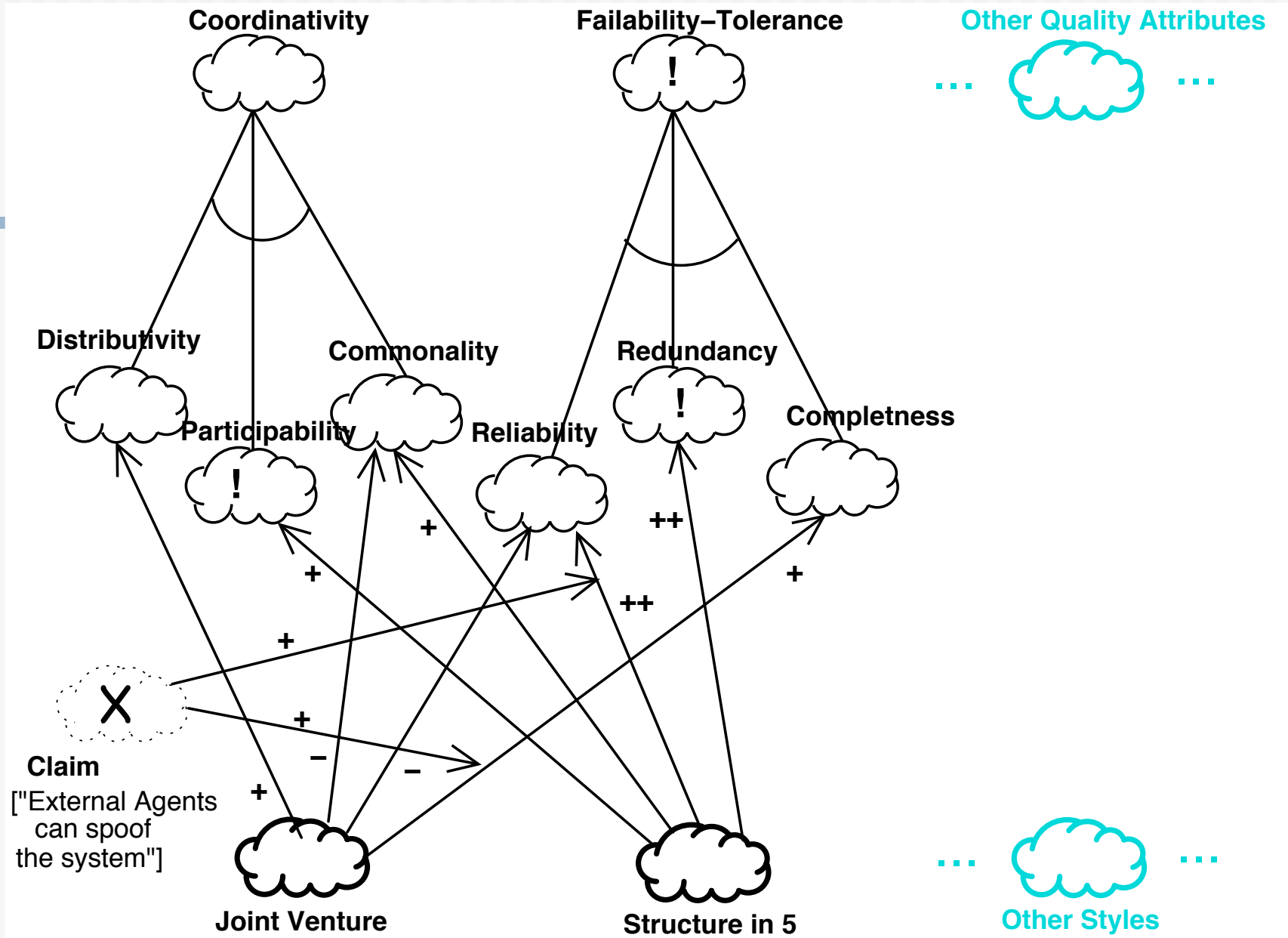


Actor diagram: *Late requirements analysis*



Reasoning about goal models cont.

- Architectural design
 - Allow to decide among different architectures
 - Find and solve possible conflicting situation among subcomponents
 - Important when sub-component(/actors)use the same resources
 - Evaluate sub-part of the design (step-by-step evaluation)
- Pattern selection
 - Patterns can be evaluated and selected with respect to
 - Their impact on goals/softgoals
 - Their impact on other patterns



Reference

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- P. Giorgini, J. Mylopoulos, and R. Sebastiani. Goal-Oriented Requirements Analysis and Reasoning in the Tropos Methodology. Engineering Applications of Artificial Intelligence, Elsevier, Volume 18/2, March 2005.
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